

Knowledge Base Partitioning

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FORMAL REASONING GROUP

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Thanks to:

- Sheila A. McIlraith (Stanford/KSL)
- Vinay K. Chaudhri (SRI)

Motivating Problem

- Exploiting structure inherent in a set of logical axioms for efficient reasoning.

Contents

1. Theorem proving with partitions
2. Automatic partitioning:
 - Algorithm SPLIT
 - Algorithm TRIANGULATE
3. Preliminary results:
 - (some of) SRI's HPKB knowledge base
 - (some of) Cyc's Spatial Axioms

A partitioning and its intersection graph

A

$\neg ok_pump \vee \neg on_pump \vee water$
 $\neg man_fill \vee water$
 $\neg man_fill \vee \neg on_pump$
 $man_fill \vee on_pump$
 $\neg water \vee \neg ok_boiler \vee \neg on_boiler$
 $\vee steam$
 $water \vee \neg steam$
 $ok_boiler \vee \neg steam$
 $on_boiler \vee \neg steam$
 $\neg steam \vee \neg coffee \vee hot_drink$
 $coffee \vee teabag$
 $\neg steam \vee \neg teabag \vee hot_drink$

A_1

(1) $\neg ok_pump \vee \neg on_pump$
 $\vee water$
 (2) $\neg man_fill \vee water$
 (3) $\neg man_fill \vee \neg on_pump$
 (4) $man_fill \vee on_pump$

A_2

$water$

(5) $\neg water \vee \neg ok_boiler \vee$
 $\neg on_boiler \vee steam$
 (6) $water \vee \neg steam$
 (7) $ok_boiler \vee \neg steam$
 (8) $on_boiler \vee \neg steam$

A_3

$steam$

(9) $\neg steam \vee \neg coffee \vee hot_drink$
 (10) $coffee \vee teabag$
 (11) $\neg steam \vee \neg teabag \vee hot_drink$

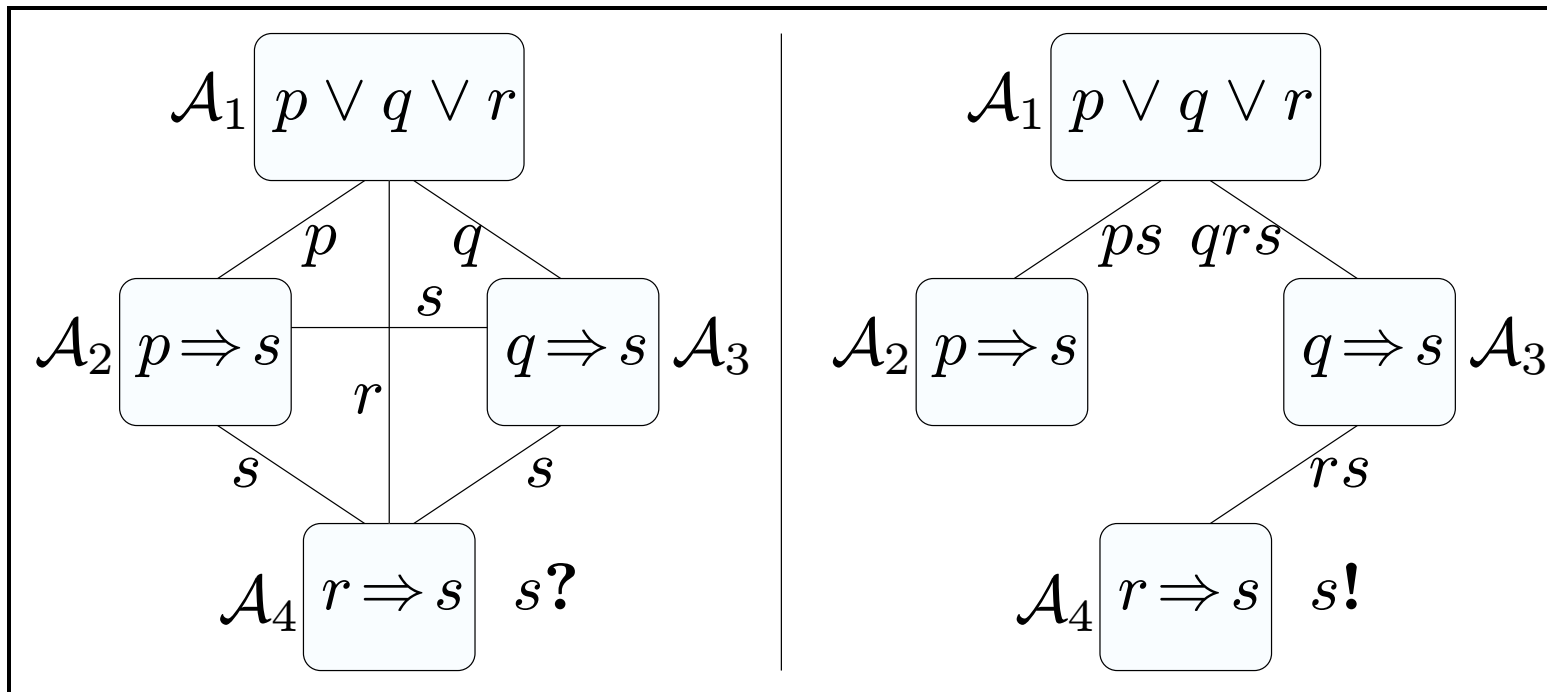
A partitioning of A and its intersection graph.

A Message-Passing Proof

Using MP to prove *hot_drink* after asserting *ok_pump* (12) in \mathcal{A}_1 and *ok_boiler* (13), *on_boiler* (14) in \mathcal{A}_2 .

\mathcal{A}_1	Resolve	Generating
(1) $\neg ok_pump \vee \neg on_pump \vee water$ (2) $\neg man_fill \vee water$ (3) $\neg man_fill \vee \neg on_pump$ (4) $man_fill \vee on_pump$	(2),(4)	$on_pump \vee water$ (m1)
	(m1),(1)	$ok_pump \vee water$ (m2)
	(m2),(12)	$water$ (m3)
		clause <i>water</i> passed from \mathcal{A}_1 to \mathcal{A}_2
\mathcal{A}_2 <div style="text-align: center; margin-bottom: 5px;">$water$</div> (5) $\neg water \vee \neg ok_boiler \vee \neg on_boiler \vee steam$ (6) $water \vee \neg steam$ (7) $ok_boiler \vee \neg steam$ (8) $on_boiler \vee \neg steam$	(m3),(5)	$ok_boiler \wedge on_boiler \Rightarrow steam$ (m4)
	(m4),(13)	$\neg on_boiler \vee steam$ (m5)
	(m5),(14)	$steam$ (m6)
		clause <i>steam</i> passed from \mathcal{A}_2 to \mathcal{A}_3
\mathcal{A}_3 <div style="text-align: center; margin-bottom: 5px;">$steam$</div> (9) $\neg steam \vee \neg coffee \vee hot_drink$ (10) $coffee \vee teabag$ (11) $\neg steam \vee \neg teabag \vee hot_drink$	(9),(10)	$\neg steam \vee teabag \vee hot_drink$ (m7)
	(m7),(11)	$\neg steam \vee hot_drink$ (m8)
	(m8),(m6)	hot_drink (m9)

Graphs with Cycles



An intersection graph before (left) and after (right) applying **BREAK-CYCLES**.

Soundness and Completeness of MP

Craig's Interpolation Theorem: *If $\alpha \vdash \beta$, then there is a formula γ involving only symbols common to both α and β , such that $\alpha \vdash \gamma$ and $\gamma \vdash \beta$.*

Theorem 1 *Message-Passing with Break-Cycles is sound and complete, if the reasoning procedure in each partition is complete for consequence finding.*

(e.g., $\alpha = \mathcal{A}_1, \beta = \mathcal{A}_2 \Rightarrow Q$)

A Good partitioning

(Amir & McIlraith 2000): Minimize

1. The total number of symbols contained in all links to/from node i .
2. The number of axioms in each partition.
3. p - the number of partitions.

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Optimal partitioning

1. Finding optimal partitioning is NP-hard (Arnborg et al. 1987).
2. Can the optimal be approximated with constant factor in polynomial time? – Open question
3. Many partitioning algorithms so far:
(Lagergren and Arnborg 1991), (Reed 1992),
(Kloks 1994), (Bodlaender & Kloks 1991),
(Bodlaender 1996), (Becker & Geiger 1996),
(Shoikhet & Geiger 1997).

Automated Partitioning using SPLIT

A

$\neg ok_pump \vee \neg on_pump \vee water$

$\neg man_fill \vee \neg on_pump$

$\neg water \vee \neg ok_boiler \vee \neg on_boiler \vee steam$

$ok_boiler \vee \neg steam$

$\neg steam \vee \neg coffee \vee hot_drink$

$\neg steam \vee \neg teabag \vee hot_drink$

$\neg man_fill \vee water$

$man_fill \vee on_pump$

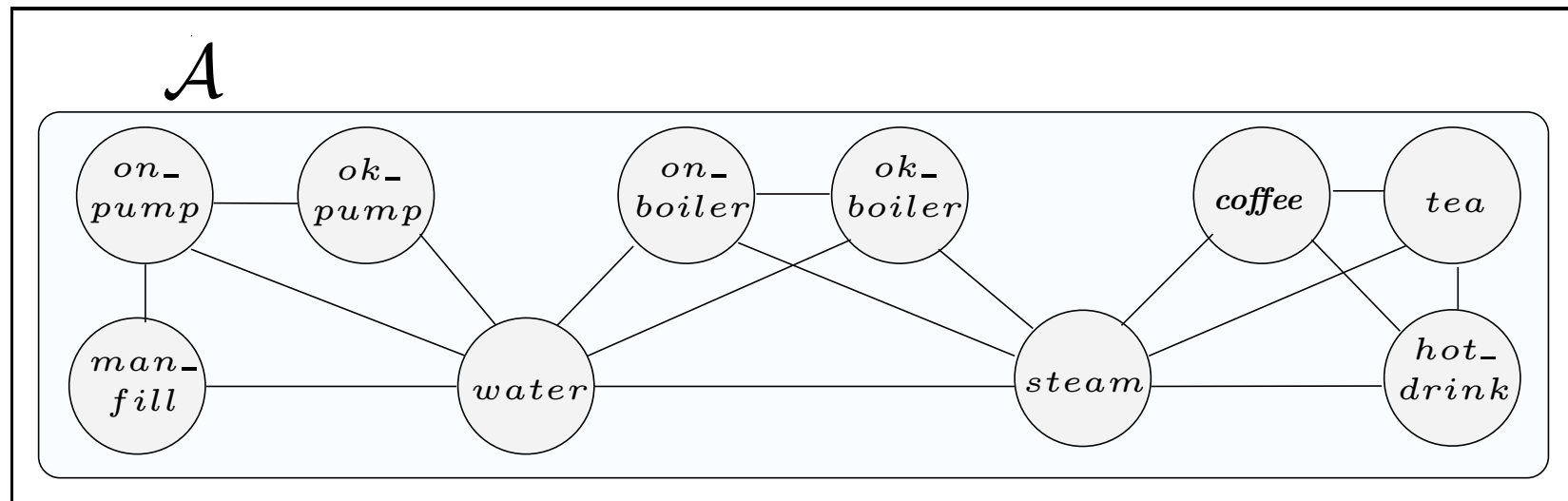
$water \vee \neg steam$

$on_boiler \vee \neg steam$

$coffee \vee teabag$

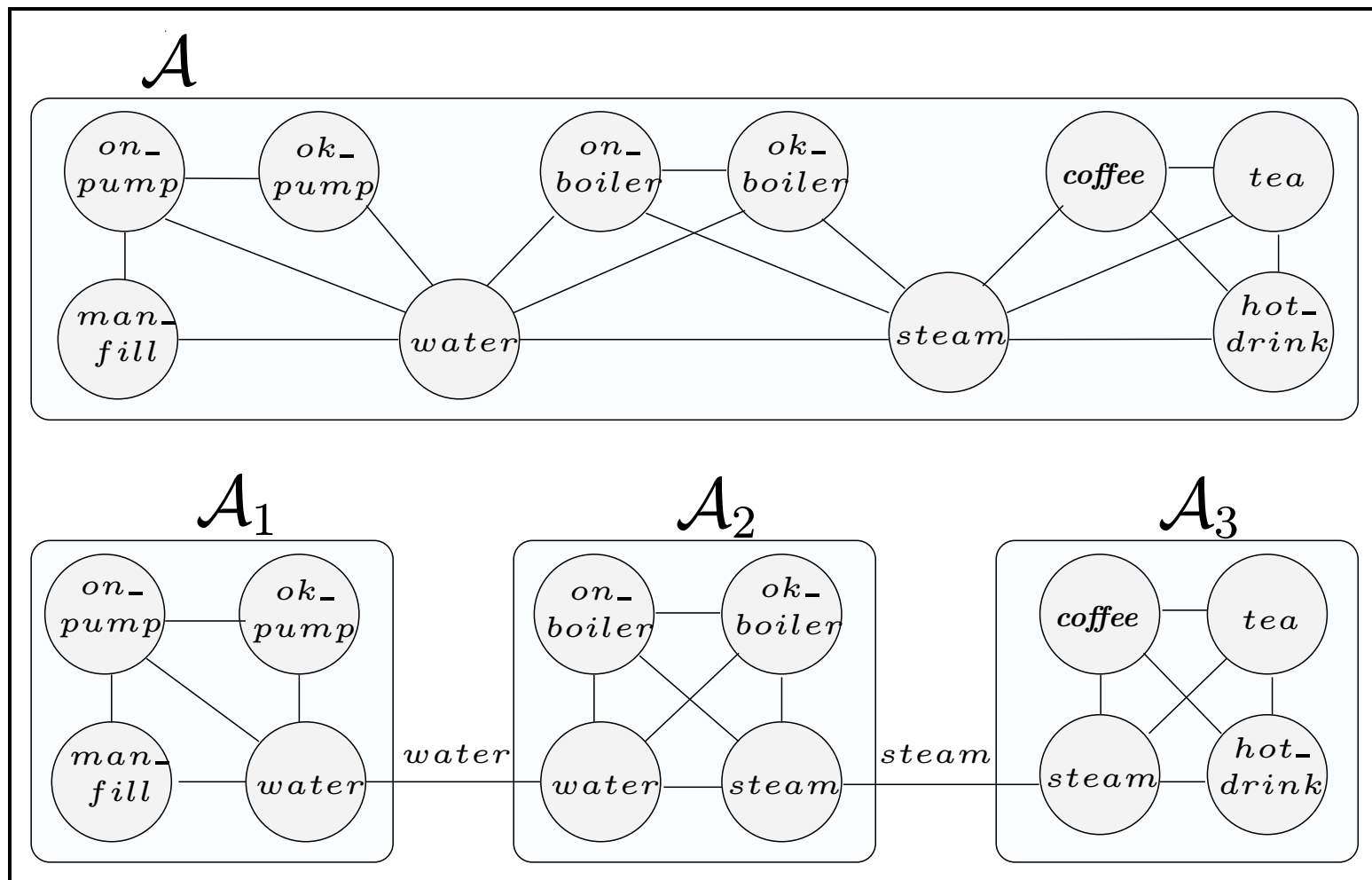
A's theory

Automated Partitioning using SPLIT



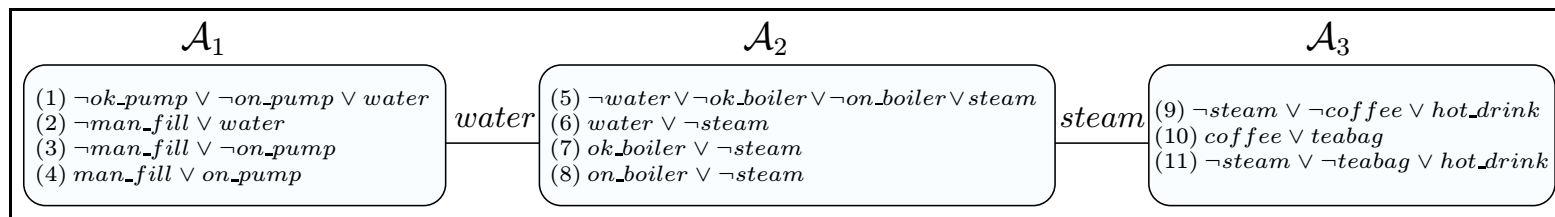
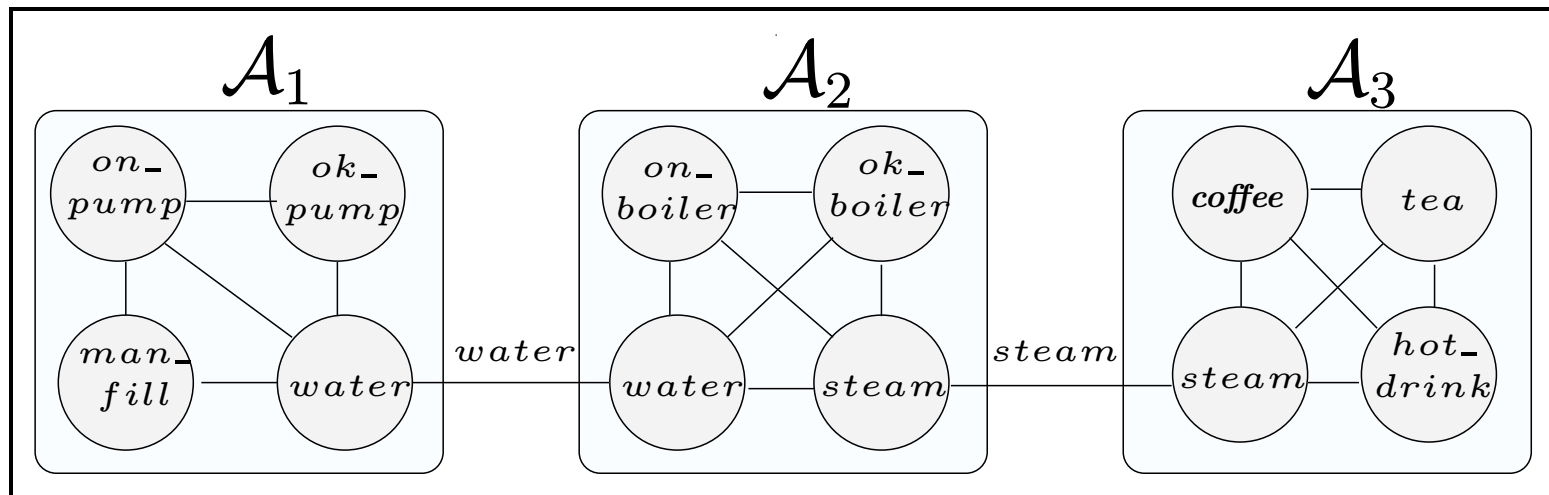
A's symbols graph

Automated Partitioning using SPLIT



Decomposing *A*'s symbols graph with *vertex min-cut*

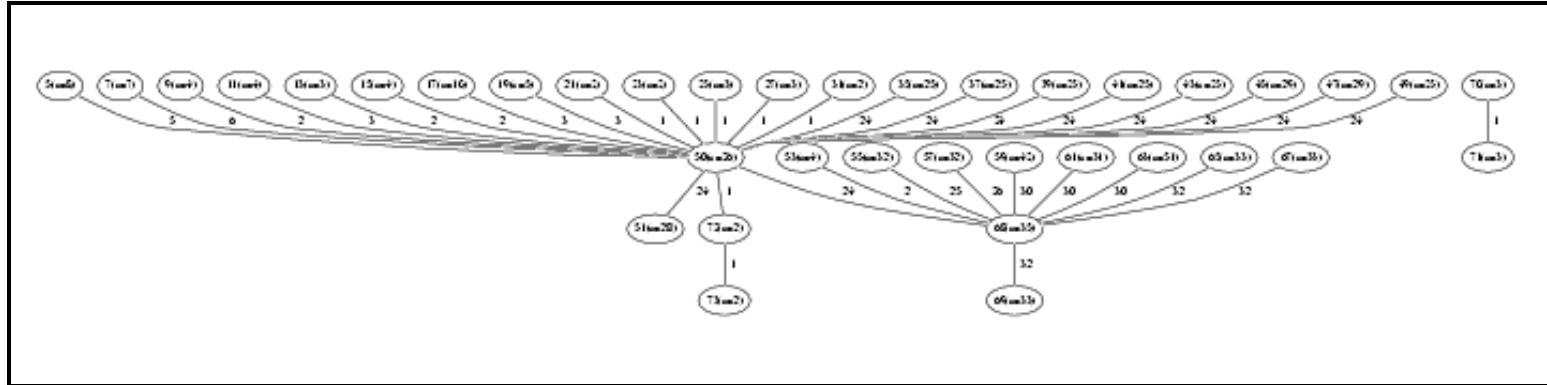
Automated Partitioning using SPLIT



Analysis of SPLIT

- At each iteration, SPLIT computes a minimum vertex separator.
- SPLIT is a greedy, recursive algorithm.
- It takes time $O(|V|^{\frac{5}{2}} * |E|)$: linear dependency on the number of clauses, and small polynomial dependency on the number of propositional symbols.
- The algorithm should stop only when the subgraph is a clique (assuming $f_{SAT}(m) = O(2^{\alpha m})$).

Using SPLIT on (some of) Cyc's Spatial Axioms



223 axioms over 142 symbols.

Time for finding the decomposition: 5 seconds.

Largest partition: 51 symbols

Largest link: 32 symbols

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Automated Partitioning using TRIANG.

A

$\neg ok_pump \vee \neg on_pump \vee water$

$\neg man_fill \vee \neg on_pump$

$\neg water \vee \neg ok_boiler \vee \neg on_boiler \vee steam$

$ok_boiler \vee \neg steam$

$\neg steam \vee \neg coffee \vee hot_drink$

$\neg steam \vee \neg teabag \vee hot_drink$

$\neg man_fill \vee water$

$man_fill \vee on_pump$

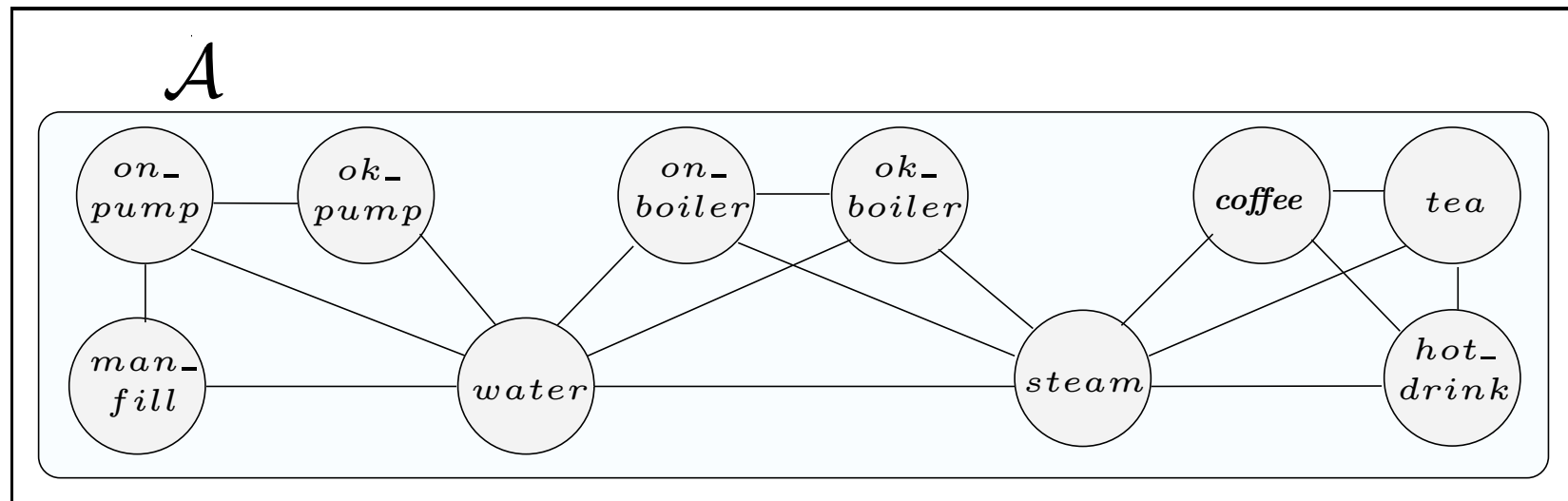
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$coffee \vee teabag$

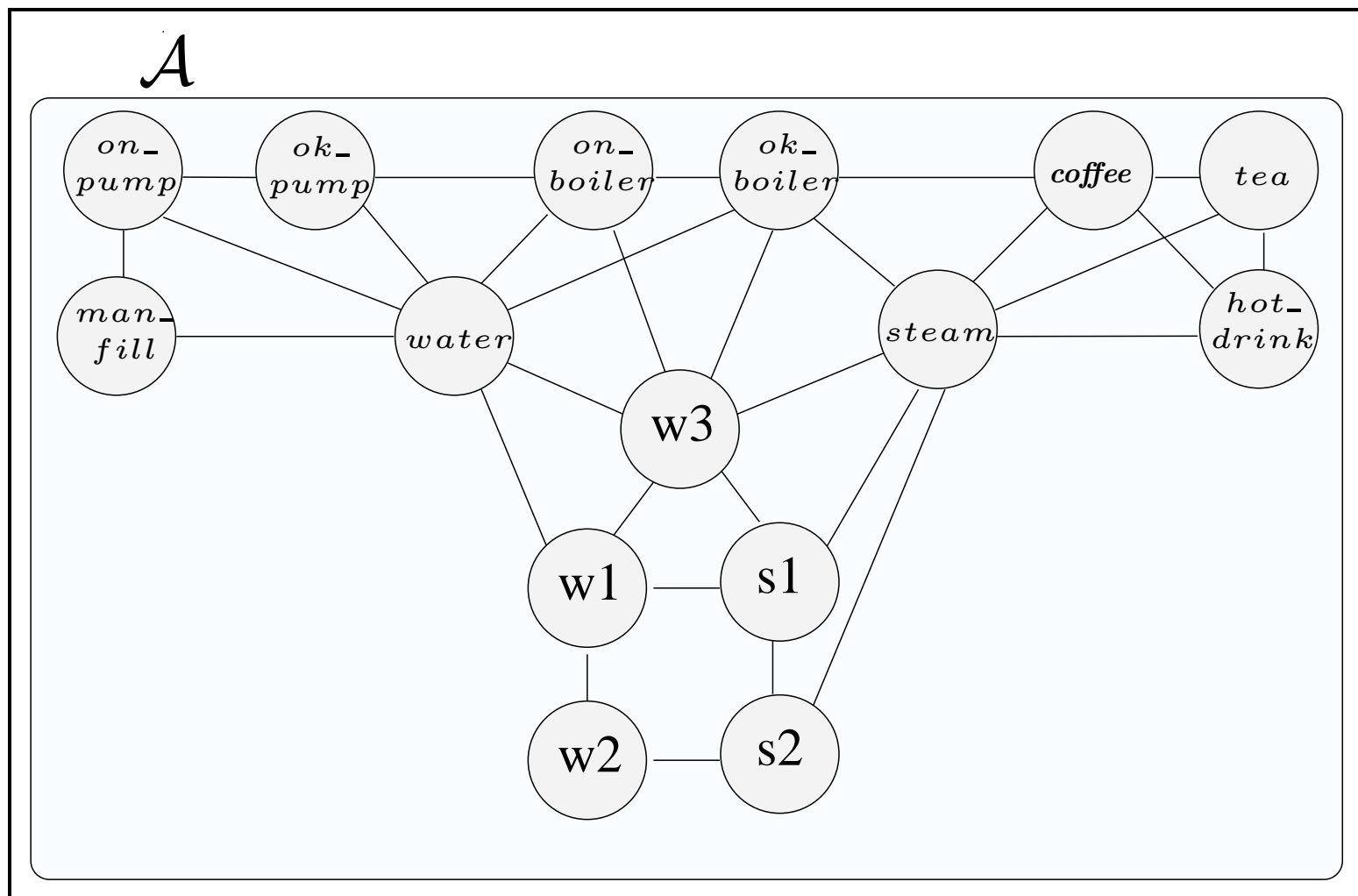
A's theory

Automated Partitioning using TRIANG.



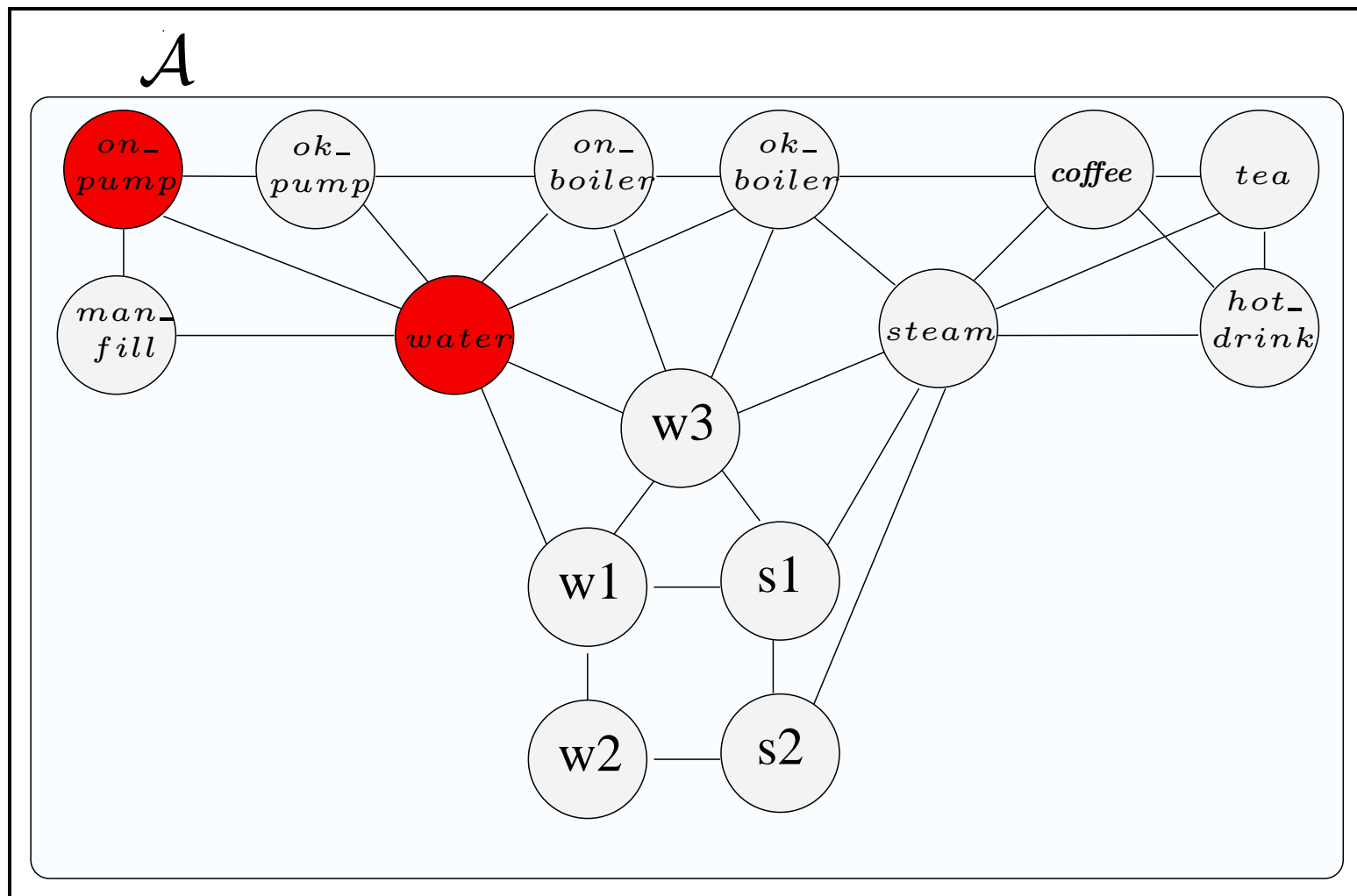
A's symbols graph

Automated Partitioning using TRIANG.



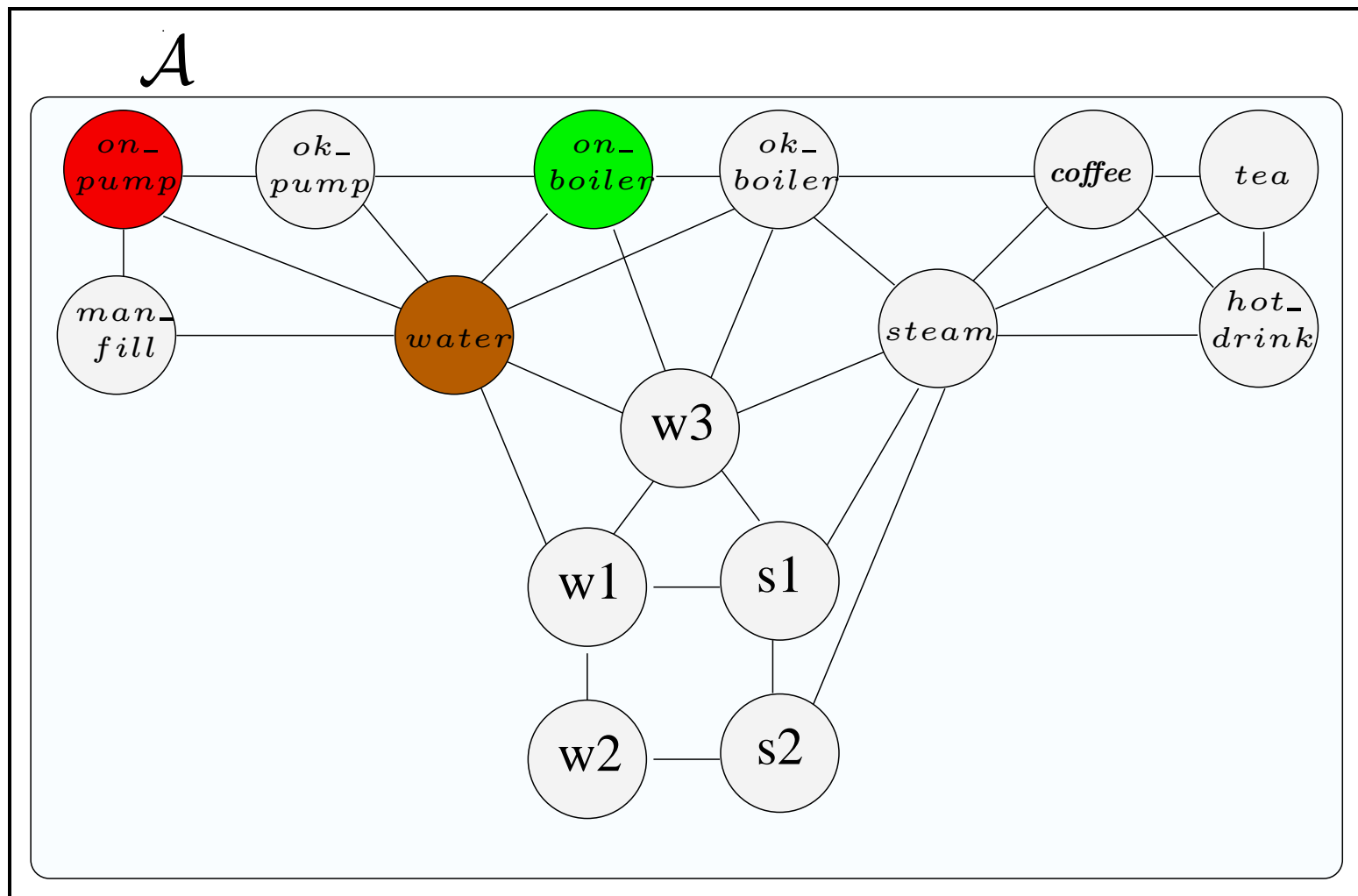
A's symbols graph (axioms/symbols appended).

Automated Partitioning using TRIANG.



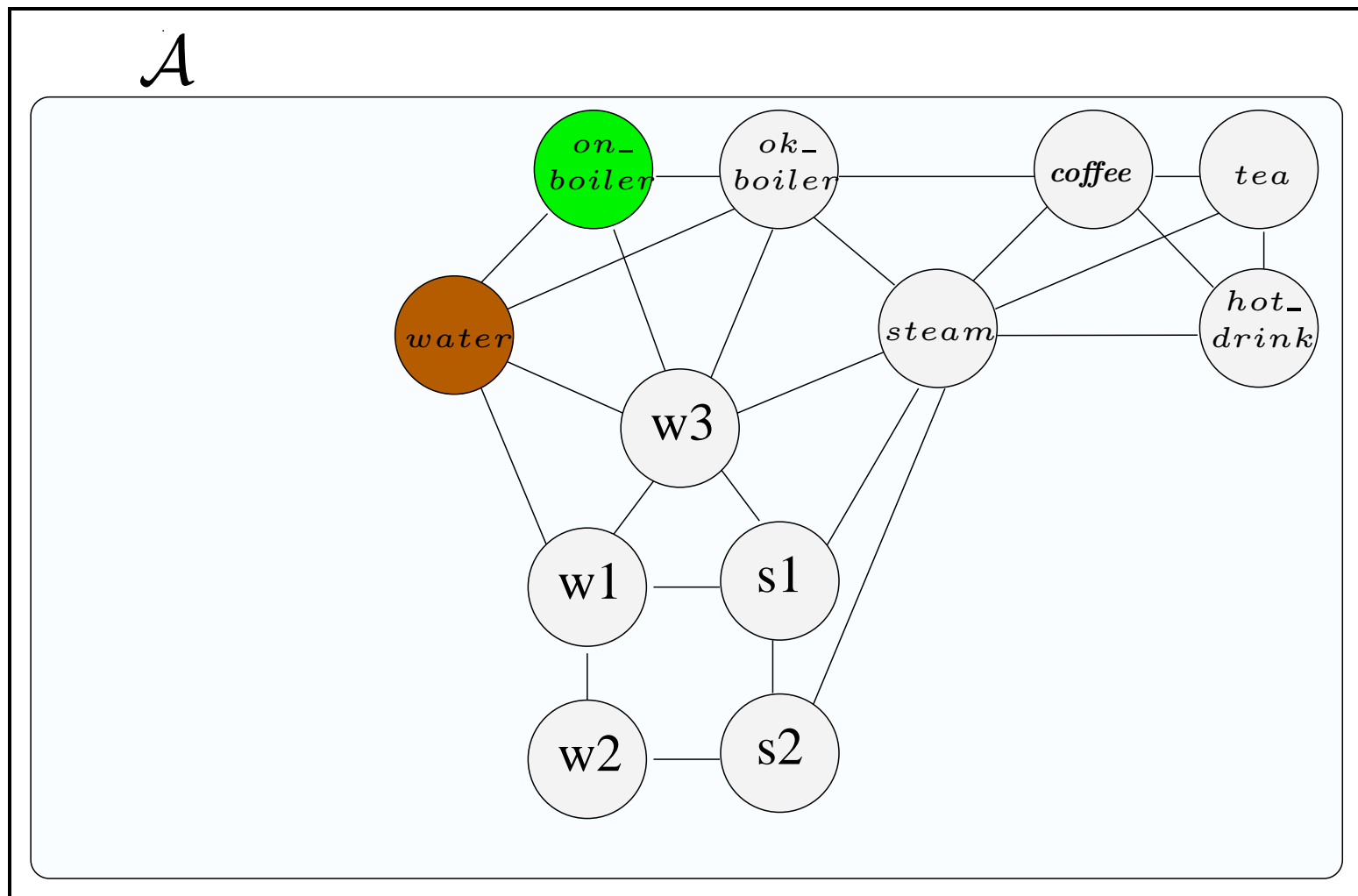
Using a *balanced separator* of previous separators

Automated Partitioning using TRIANG.



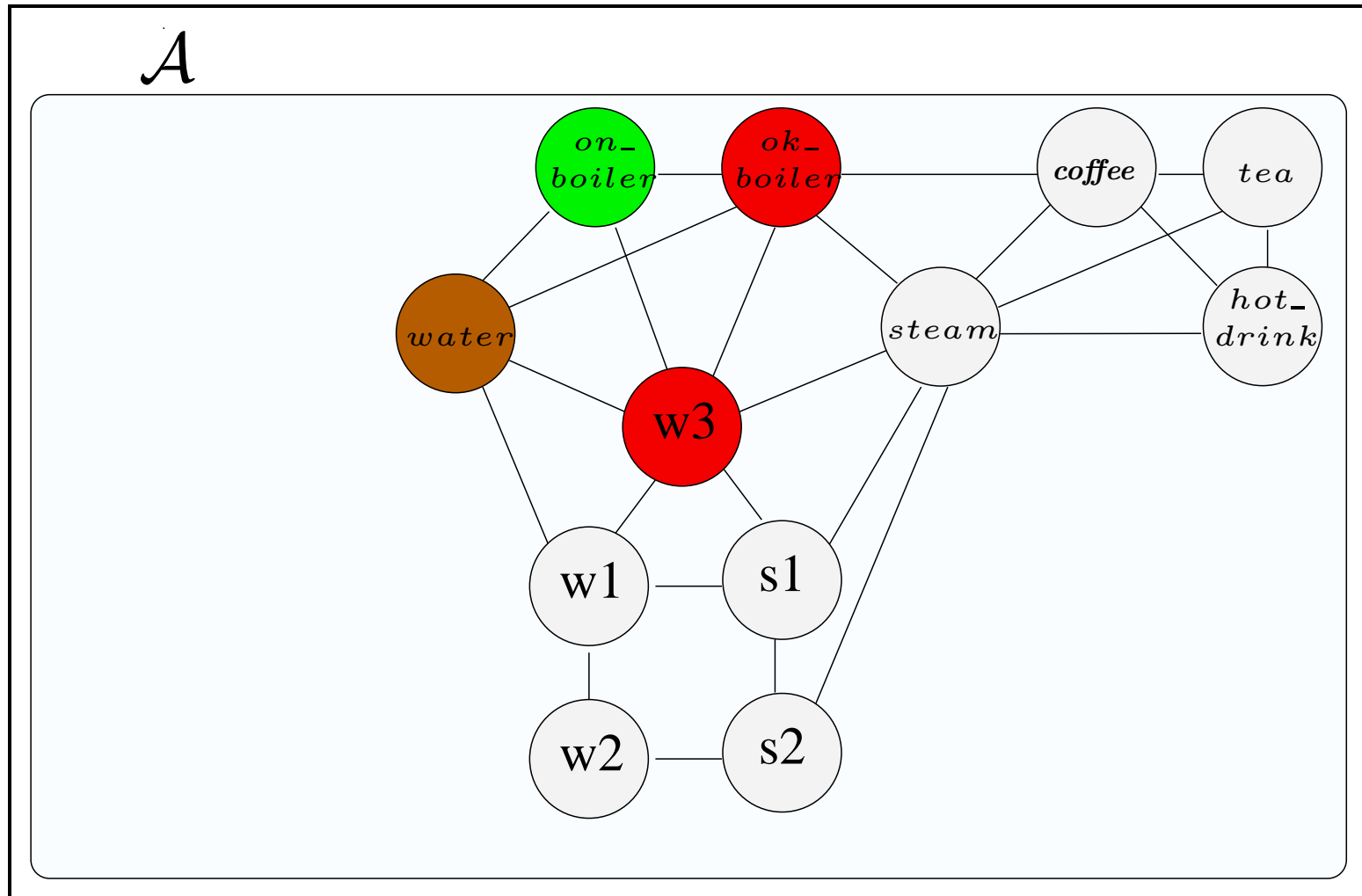
Using a *balanced separator* of previous separators

Automated Partitioning using TRIANG.



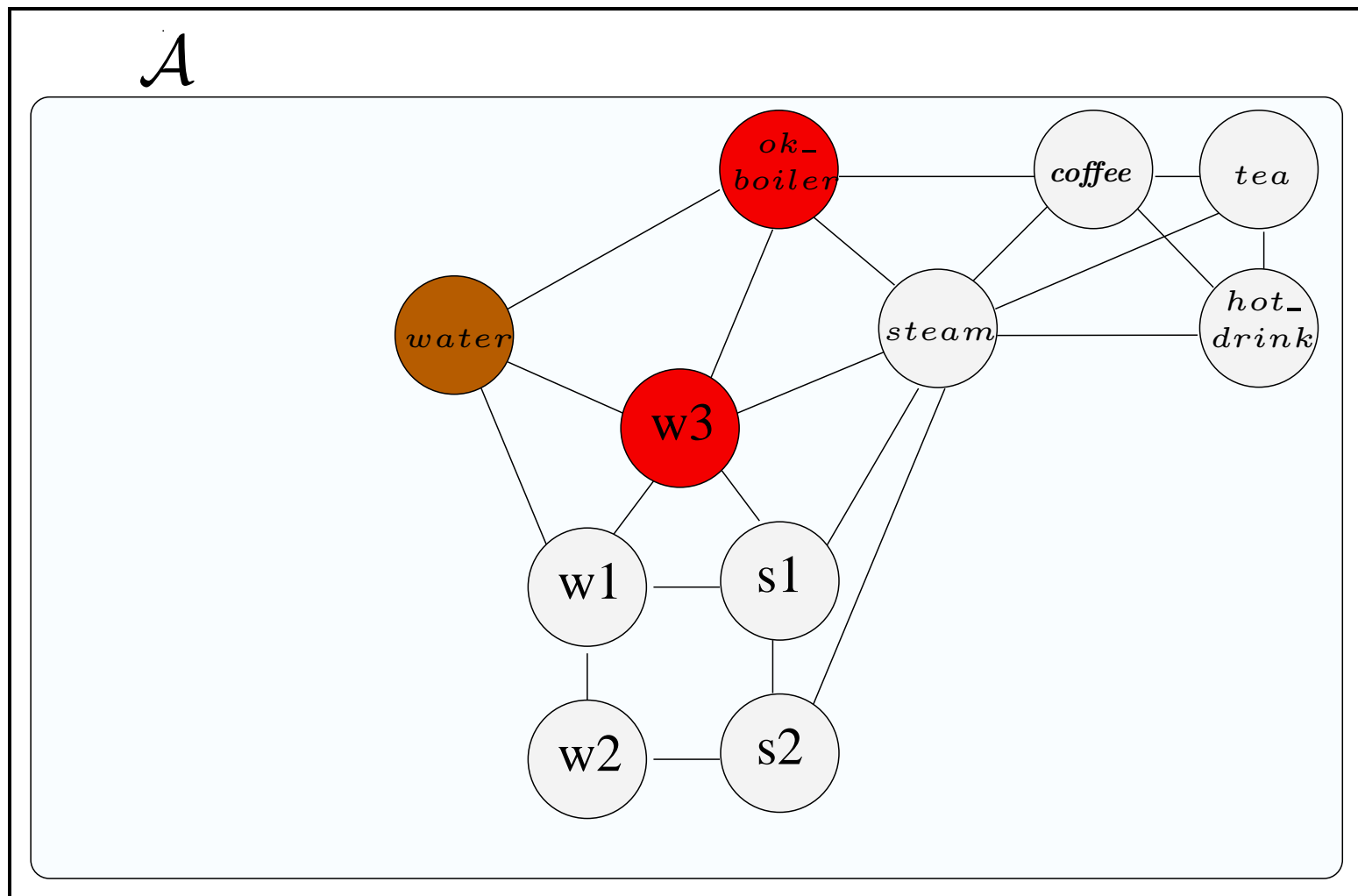
Using a *balanced separator* of previous separators

Automated Partitioning using TRIANG.



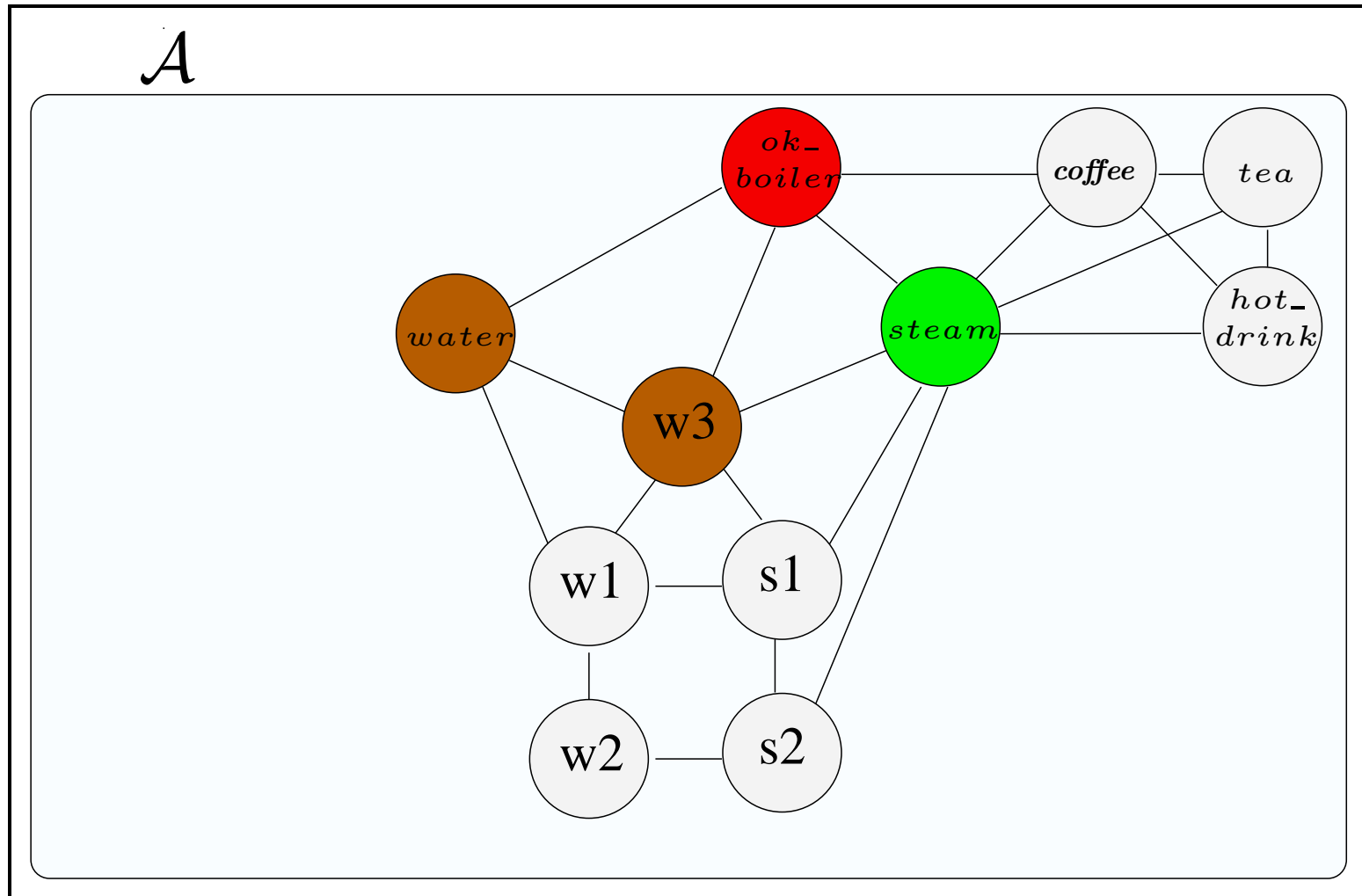
Using a *balanced separator* of previous separators

Automated Partitioning using TRIANG.



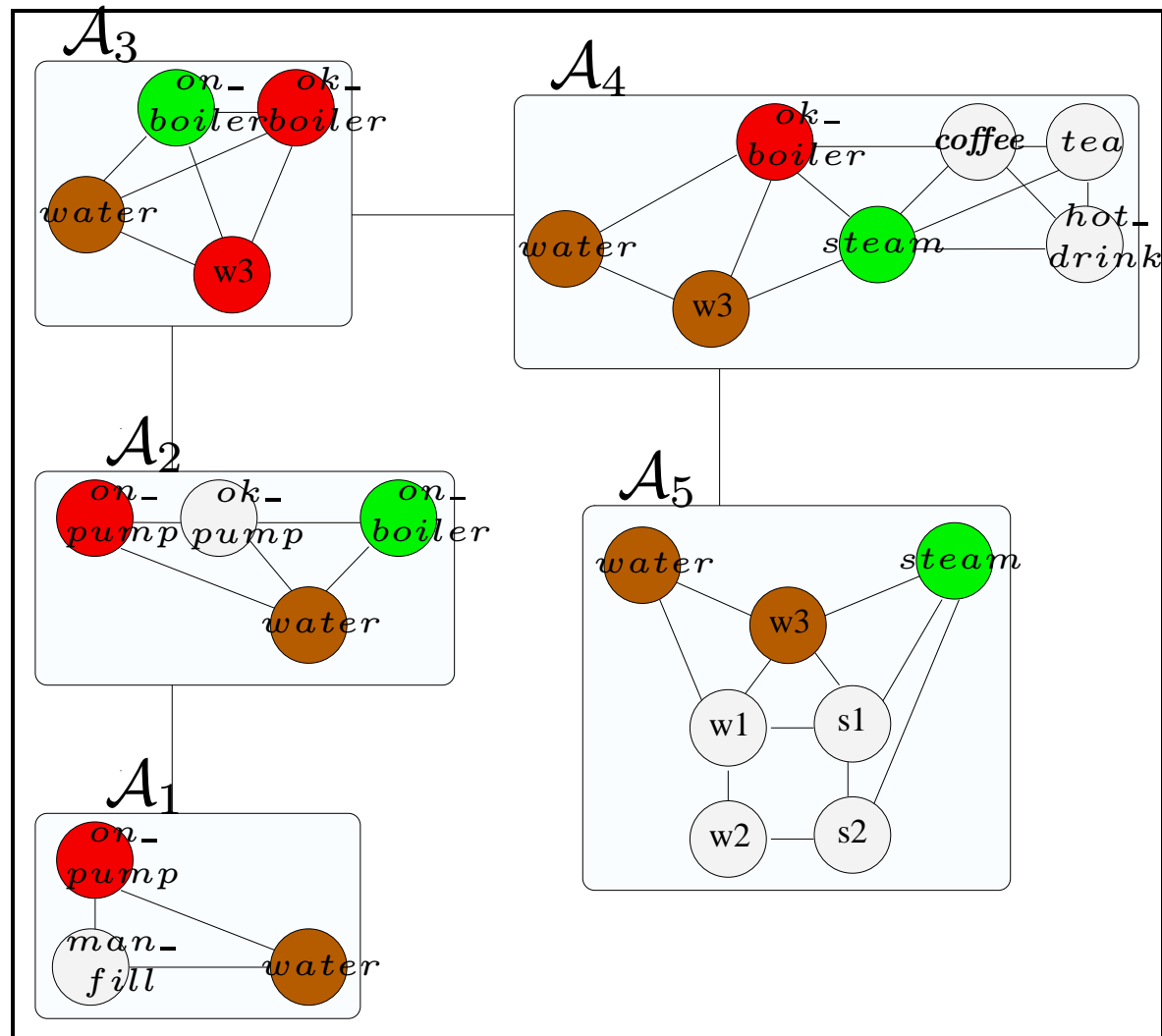
Using a *balanced separator* of previous separators

Automated Partitioning using TRIANG.



Using a *balanced separator* of previous separators

Automated Partitioning using TRIANG.



Overall decomposition graph.

Analysis of TRIANGULATE

- TRIANGULATE is guaranteed to find a factor-3 approximation to the optimal.
- It takes time $O(2^{2k} k^{\frac{5}{2}} |V|^2)$ (k is the size of the largest partition in the optimal decomposition).

Using TRIANGULATE on (some of) Cyc's Spatial Axioms

223 axioms over 142 symbols.

Time: 3.5 minutes

Largest partition: 21 symbols.

Largest link: 20 symbols.

Using TRIANGULATE on (some of) SRI's HPKB Axioms

Original Axioms

688 axioms with 570 symbols.

Time = 13 hours;

Largest partition with 70 symbols.

CNF axioms

1199 axioms with 446 symbols.

Time = 2.5 hours;

Largest partition with 58 symbols.

Summary

Motivating Problems:

- Exploiting structure inherent in a set of logical axioms for efficient reasoning.

Contribution:

- Showed that commonsense knowledge bases are decomposable.
- We provide decomposition algorithms for FOL theories.
- Good Decomposition of theories provide computational advantages for theorem provers.

Current & Future Work

- Examine other decomposition algorithms.
- Analyze the decompositions of the HPKB axioms and Cyc's axioms and compare them to human-made decompositions.
- Experimental analysis of partition-based reasoning.
- Open directions: equality, types, logics, planning, prime-implicates.

<http://www-formal.stanford.edu>

Extensions

- Relaxing completeness for consequence finding: \mathcal{L} -generation completeness
- Exploiting polarity using Lyndon's interpolation theorem
- Queries from multiple partitions
- Special treatment for first-order resolution.

<http://www-formal.stanford.edu>

Related Work

- **Bayes Nets:** e.g., (Pearl 1988), (Becker & Geiger 1996).
- **CSPs:** e.g., (Dechter & Pearl 1988), (Gottlob et al 1999).
- **Join Trees:** e.g., (Shoikhet & Geiger 1997), (Kloks 1994).
- **SAT:** e.g., (Selman & Kautz 1993), (Dechter & Rish 1994), (Park & VanGelder 1996).
- **Parallel Thm Proving:** e.g., (Bonacina 1994), (Cowen & Wyatt 1993), (Suttner 1997).
- **Combining Logical Systems:** e.g., (Nelsson & Oppen 1979), (Shostak 1984), (Baader & Schulz

1992).

- **Partitioned Representations:** e.g., (McCarthy & Buvac 1998), (Giunchiglia & Ghidini 1998), (Amir 2000).

Effectiveness of MP

Currently available measures for effectiveness of proof procedures:

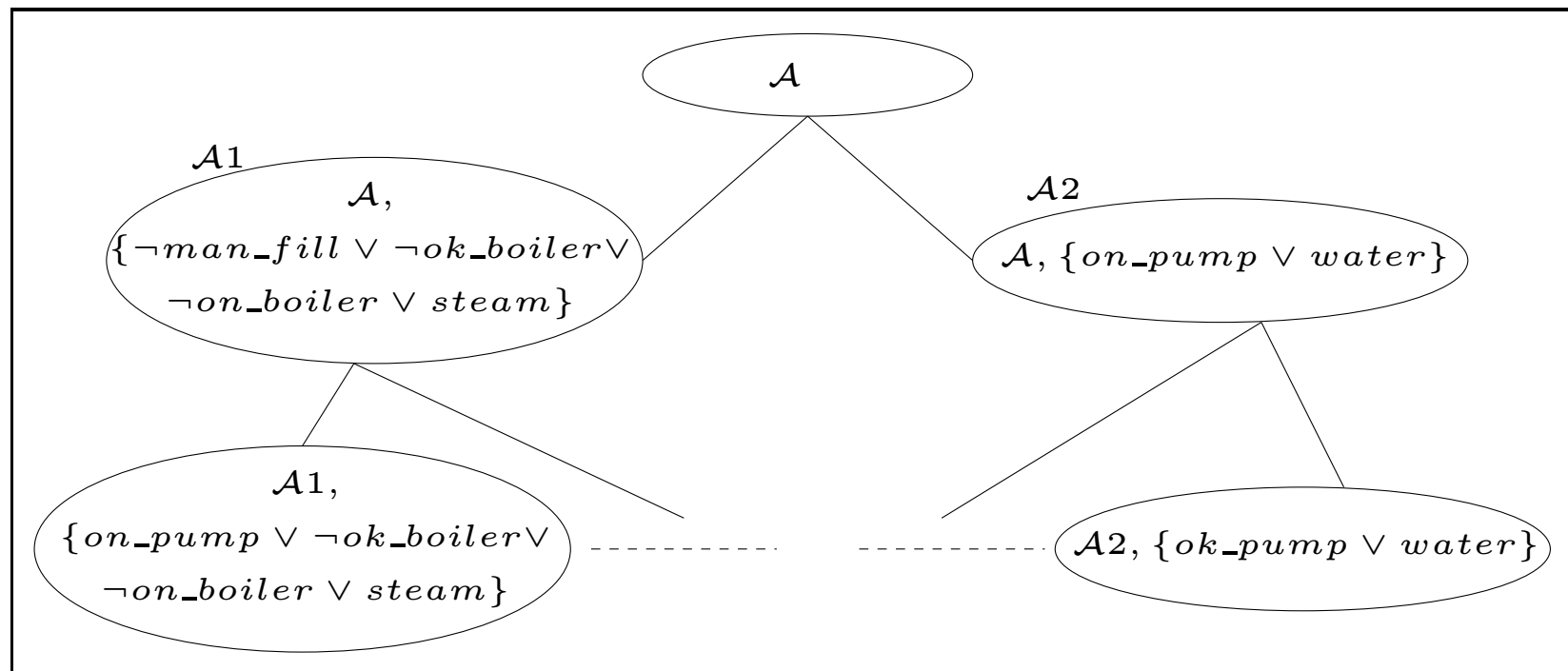
- Search-space size
- Proof length

Theorem 2 Resolution [w/restriction strategies] applied using MP has:

- *restricted search space, depending on partitions sizes and links sizes.*
- *small interpolants and short proofs, depending on links sizes.*

Effectiveness of MP: Search Space

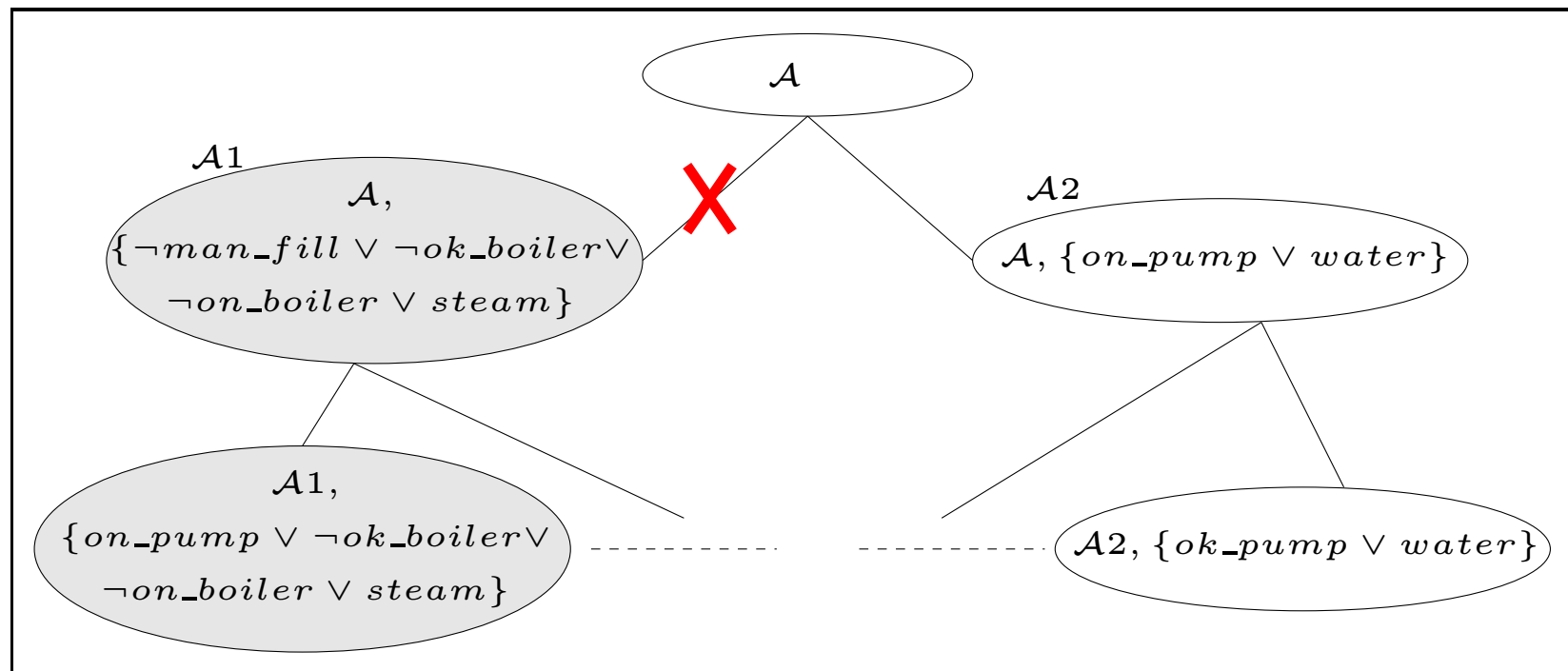
- In resolution-based inference, the search space using MP is restricted.



Linear Resolution

Effectiveness of MP: Search Space

- In resolution-based inference, the search space using MP is restricted.



Linear Resolution restricted with MP

Effectiveness of MP: Proof Length

Propositional proof lengths with/without MP:

- **Longest proof:** The longest non-MP proof is as long or longer than the longest MP proof.
- **Shortest proof:**
 - The shortest non-MP proof can be shorter than the shortest MP proof.
 - **Evidence:** The difference in length between the shortest MP proof and shortest non-MP proof is polynomially bounded, if the size of links is fixed.
 - **Evidence:** The difference is exponentially

bounded in the size of the links.